



# *XcalableMP Ver. 1.0*

Highly Productive Parallel Programming  
Language

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# Background



- MPI is widely used as a parallel programming model on distributed memory systems
  - Time-consuming
  - Complicated process
- Another programming model is needed !!
  - High performance
  - Easy to program (High productivity)



 **Development of XcalableMP (XMP)**

# e-Science Project



- XMP Working Group designs XMP specification
  - XMP Working Group consists of members from
    - academia : U. Tsukuba, U. Tokyo, Kyoto U. and Kyusyu U.
    - research labs : RIKEN AICS, NIFS, JAXA, JAMSTEC/ES
    - industries : Fujitsu, NEC, Hitachi
  - Specification Version 1.0 is released !! (Nov. 2011)
- University of Tsukuba develops an Omni XMP compiler as a reference implementation
  - the K Computer, CRAY platforms(HECToR), Linux cluster
- Evaluation of Performance and Productivity



**Directive-based language extension for  
Scalable and performance-aware Parallel Programming**

[\[Japanese\]](#)

[\[Home\]](#) [\[Download\]](#) [\[Overview\]](#) [\[Manual\]](#) [\[Tutorial\]](#) [\[Publication\]](#) [\[Support\]](#)

## What's XcalableMP

XcalableMP is a directive-based language extension which allows users to develop parallel programs for distributed memory systems easily and to tune the performance by having minimal and simple notations. The specification is being designed by XcalableMP Specification Working Group which consists of members from academia and research labs to industries in Japan.

- [Specification of XcalableMP, version 1.0](#) (Nov. 14. 2011)

If you have any comments and requests, please contact to [Prof. Mitsuhsa Sato](#)(msato at cs.tsukuba.ac.jp). Your comments and contributions will be appreciated !!

The features of XcalableMP are summarized as follows:

- XcalableMP supports typical parallelization based on the data parallel paradigm and work mapping under "global view programming model", and enables parallelizing the original sequential code using minimal modification with simple directives, like OpenMP. Many ideas on "global-view" programming are inherited from HPF (High Performance Fortran).
- The important design principle of XcalableMP is "performance-awareness". All actions of communication and synchronization are taken by directives, different from automatic parallelizing compilers. The user should be aware of what happens by XcalableMP directives in the execution model on the distributed memory architecture.
- XcalableMP also includes a CAF-like PGAS (Partitioned Global Address Space) feature as "local-view" programming.
- Extension of existing base languages with directives is useful to reduce rewriting and educational costs. XcalableMP APIs are

<http://www.xcalablemp.org/>

WPSE2012@Kobe, Japan

# Agenda

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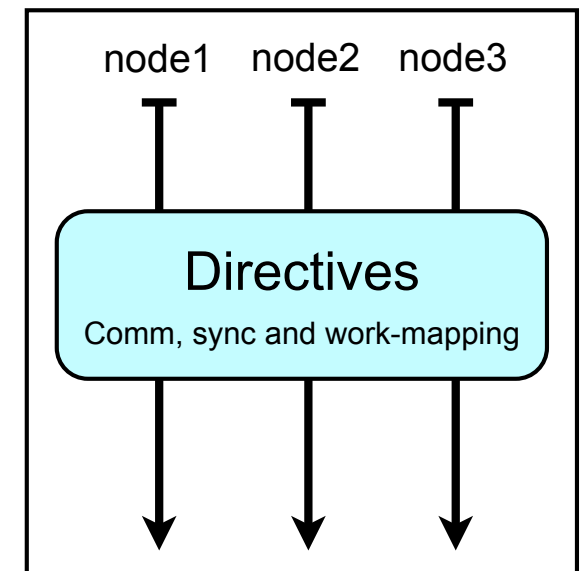


- Overview of XMP
- XMP Programming Model
- Evaluation of Performance and Productivity of XMP

# Overview of XMP



- XMP is a directive-based language extension like OpenMP and HPF based on C and Fortran95
  - To reduce code-writing and educational costs
- The basic execution model of XMP is SPMD
  - A thread starts execution in each node independently (as in MPI)
- “Performance-awareness” for explicit communication, sync. and work-mapping
  - All actions occur when thread encounters directives or XMP’s extended syntax
  - XMP compiler generates communication only where a user inserts them to facilitate performance tuning



# XMP Code Example



## XMP C version

```
int array[100];
#pragma xmp nodes p(*)
#pragma xmp template t(0:99)
#pragma xmp distribute t(block) onto p
#pragma xmp align array[i] with t(i)

main(){
#pragma xmp loop on t(i) reduction(+:res)
    for(i = 0; i < 100; i++){
        array[i] = func(i);
        res += array[i];
    }
}
```

**data  
distribution**

**work mapping  
& reduction**

# XMP Code Example



## XMP Fortran version

```
real a(100)
!$xmp nodes p(*)
!$xmp template t(100)
!$xmp distribute t(block) onto p
!$xmp align a(i) with t(i)
:
!$xmp loop on t(i) reduction(+:res)
  do i=1, 100
    a(i) = func(i)
    res = res + a(i)
  enddo
```

**data  
distribution**

**work mapping  
& reduction**



# The same code written in MPI

```
int array[100];

main(int argc, char **argv){
    MPI_Init(&argc, &argv);
    MPI_Comm_rank(MPI_COMM_WORLD, &rank);
    MPI_Comm_size(MPI_COMM_WORLD, &size);
    dx = 100/size;
    llimit = rank * dx;
    if(rank != (size -1)) ulimit = llimit + dx;
    else ulimit = 100;

    temp_res = 0;
    for(i=llimit; i < ulimit; i++){
        array[i] = func(i);
        temp_res += array[i];
    }

    MPI_Allreduce(&temp_res, &res, 1, MPI_INT, MPI_SUM, MPI_COMM_WORLD);
    MPI_Finalize( );
}
```

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# Programming Model

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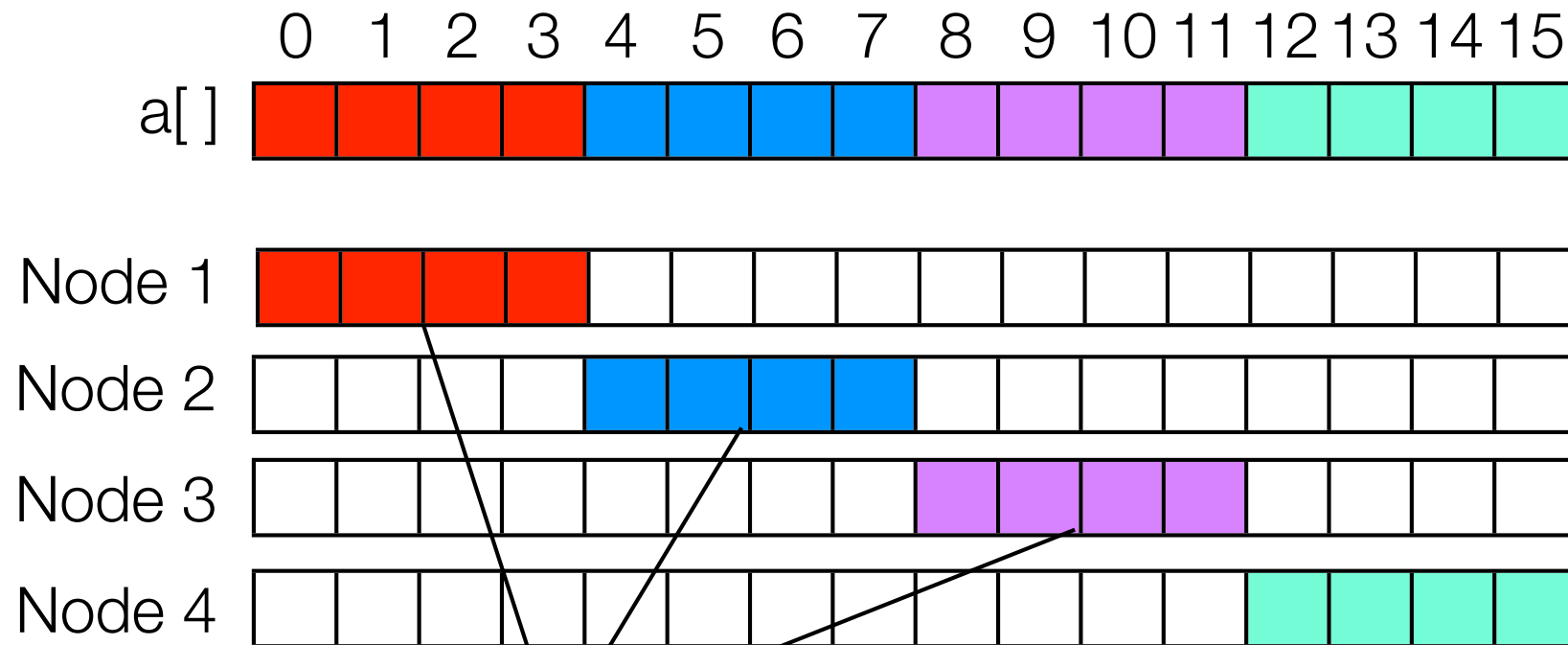
- Global View Model (like as HPF)
  - programmer describes data distribution, work mapping, communication and sync. by using directives
  - supports typical techniques for data-mapping and work-mapping
  - rich communication and sync. directives, such as “shadow”, “reflect” and “gmove”
- Local View Model (like as Coarray Fortran)
  - enables programmer to transfer data by using one-sided comm. easily

# Data Distribution



- The directives define a data distribution among nodes

```
#pragma xmp nodes p(4)  
#pragma xmp template t(0:15)  
#pragma xmp distribute t(block) on p  
#pragma xmp align a[i] with t(i)
```



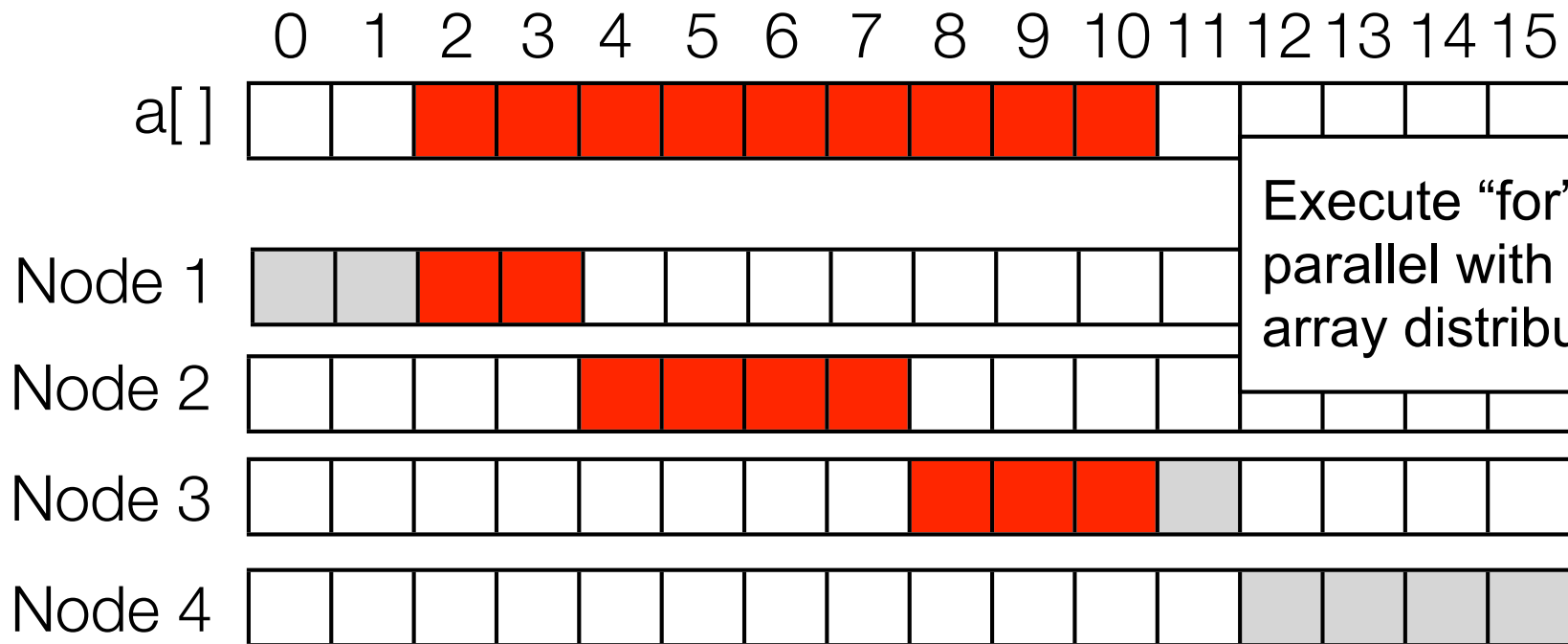
Distributed Array

# Parallel Execution of loop

- Loop directive is inserted before loop statement

```
#pragma xmp loop on t(i)  
for(i=2;i<=10;i++){...}
```

```
#pragma xmp nodes p(4)  
#pragma xmp template t(0:15)  
#pragma xmp distribute t(block) on p  
#pragma xmp align a[i] with t(i)
```



Execute "for" loop in parallel with affinity to array distribution

Each node computes **Red elements** in parallel

# Example of Data Mapping

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**block**



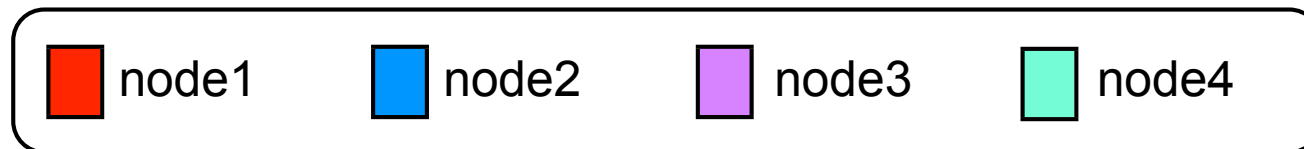
**cyclic**



**block-cyclic**  
**(block size = 3)**

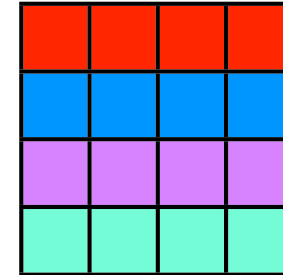


**generalized-block**

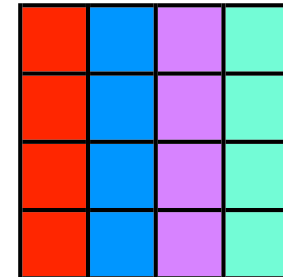


# Multi Dimensional Array

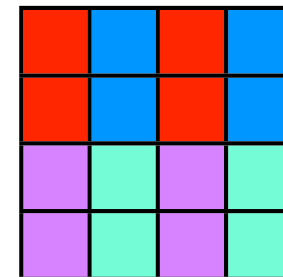
```
#pragma xmp distribute t(block) onto p  
#pragma xmp align a[i][*] with t(i)
```



```
#pragma xmp distribute t(block) onto p  
#pragma xmp align a[*][i] with t(i)
```



```
#pragma xmp distribute t(block,cyclic) onto p  
#pragma xmp align a[i][j] with t(i,j)
```



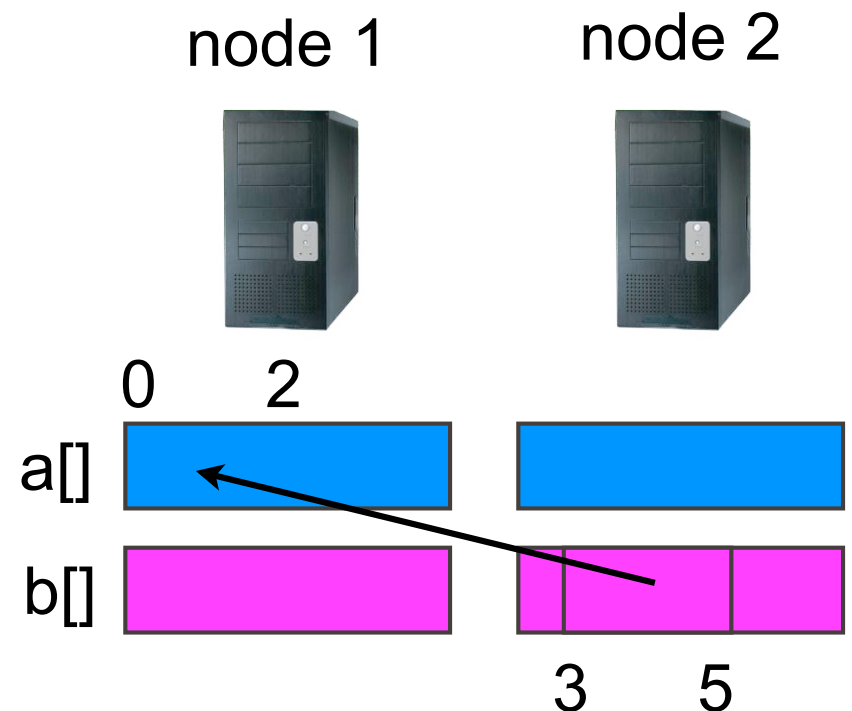
node1 node2 node3 node4

# Local View Model

- One-sided communication for local data(Put/Get)
- In XMP Fortran, this function is compatible with that of CAF
- In XMP C, C is extended to support the *array section* notation
- Uses GASNet/ARMCI, which are high-performance communication layer

```
#pragma xmp coarray b:[*]  
:  
if(me == 1)  
  a[0:3] = b[3:3]:[2]; // Get
```

base                  length                  node number





# Other directives

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<b>Directive</b>	<b>Function</b>
reduction	Aggregation
bcast	Broadcast
barrier	Synchronization
shadow/reflect	Create shadow region/sync.
gmove	Transfer for distributed data

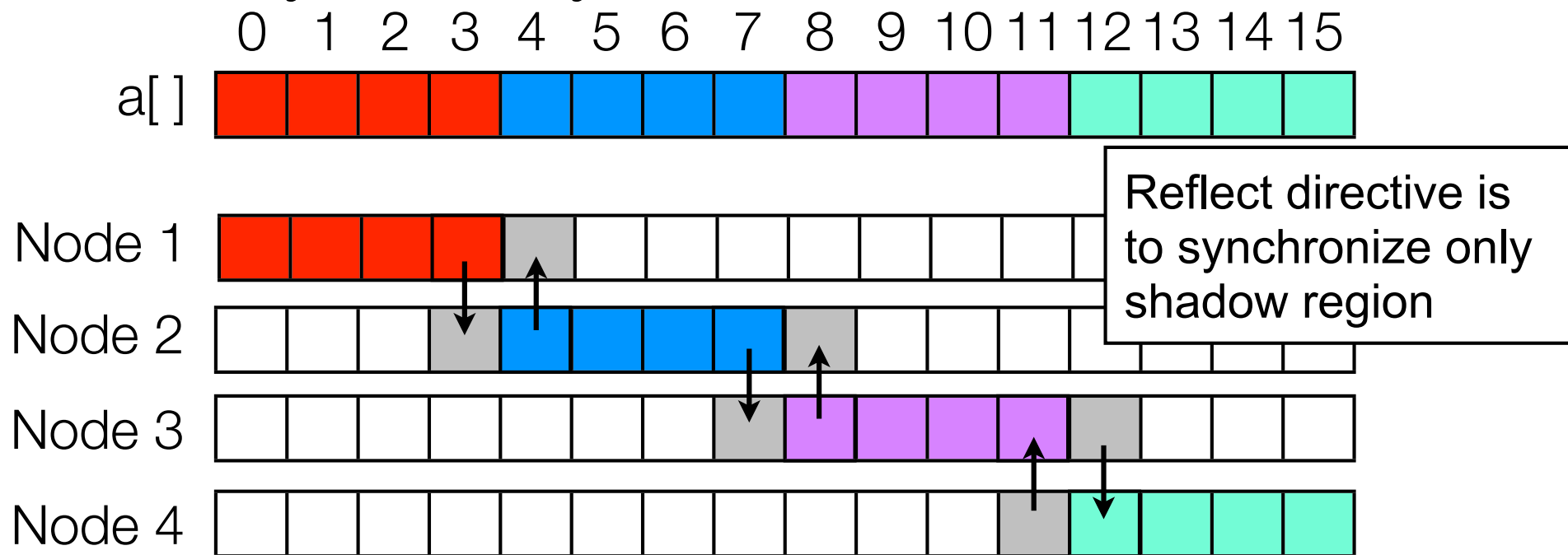
# shadow/reflect directives



- If neighbor data is required, then only shadow area can be synchronized

```
#pragma xmp shadow a[1:1]
```

- Shadow directive defines width of shadow area
- $b[i] = array[i-1] + array[i+1];$



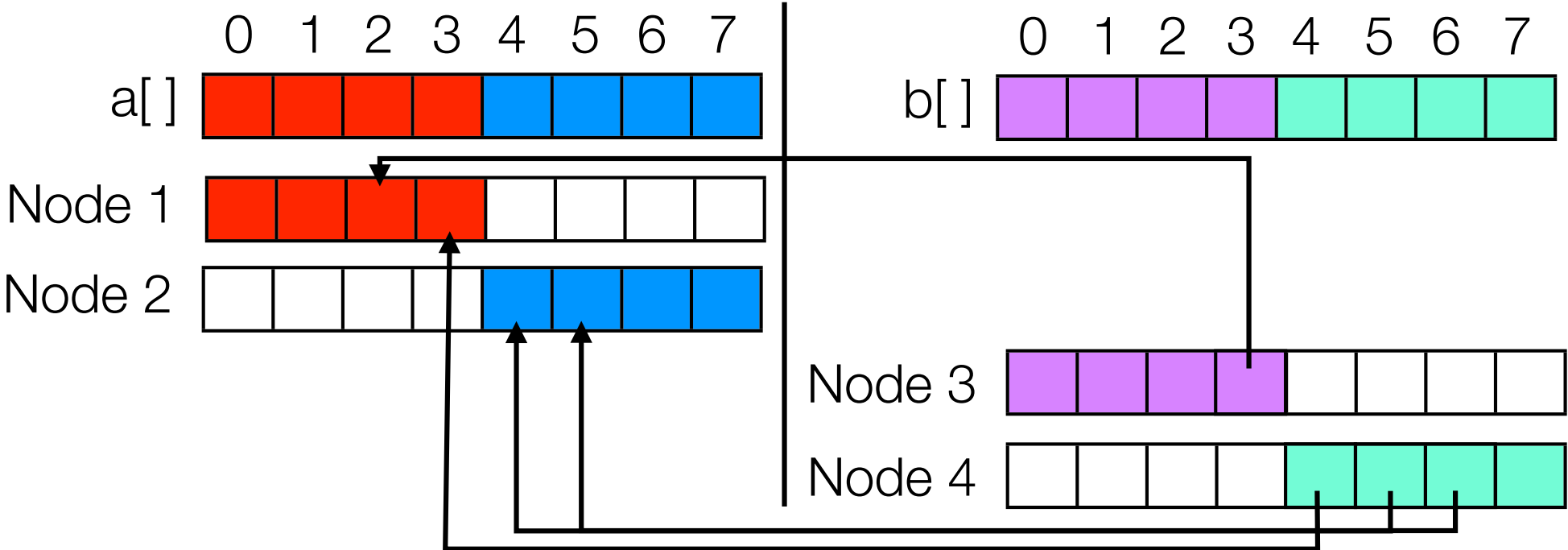
```
#pragma xmp reflect (array)
```

# Gmove directive

- Communication for distributed array
  - uses array section notation in XMP C
  - Programmer doesn't need to know where each data is distributed

```
#pragma xmp gmove  
a[2:4] = b[3:4];
```

base  
length



# Agenda

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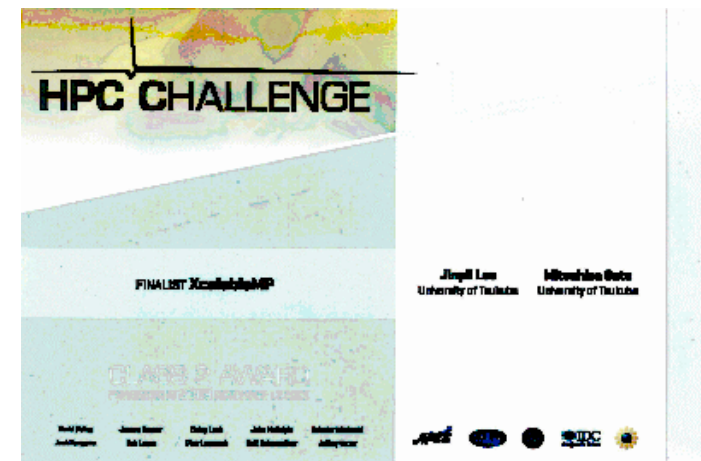


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# Evaluation



- Examines the performance and productivity of XMP
- Implementations of benchmarks
  - NAS Parallel Benchmarks
    - CG, EP, IS, BT, LU, FT, MG
  - HPC Challenge Benchmarks
    - HPL, FFT, RandomAccess, STREAM
    - Finalist of HPCC Class2 in SC10 and SC09
  - Laplace Solver
  - Himeno Benchmark, and so on



# Environment

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- T2K Tsukuba System
  - CPU : AMD Opteron Quad-Core 8356 2.3GHz (4 sockets)
  - Memory : DDR2 667MHz 32GB
  - Network : Infiniband DDR(4rails) 8GB/s



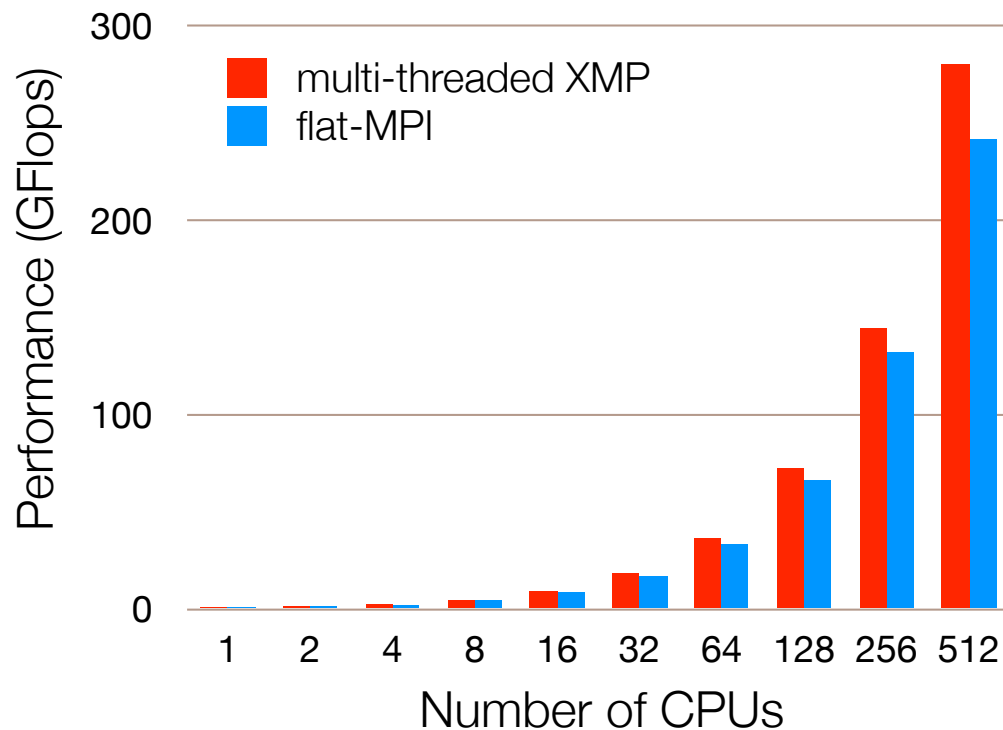
# Laplace Solver



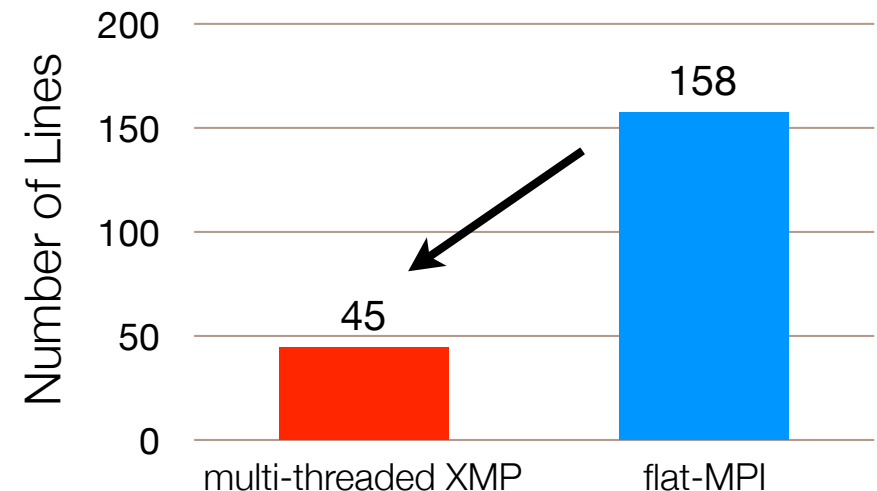
- “thread” clause for multicore cluster
- shadow/reflect directive

```
#pragma xmp loop (x, y) on t(x, y) threads  
for(y = 1; y < N-1; y++)  
  for(x = 1; x < N-1; x++)  
    tmp_a[y][x] = a[y][x];
```

## Performance



## Productivity



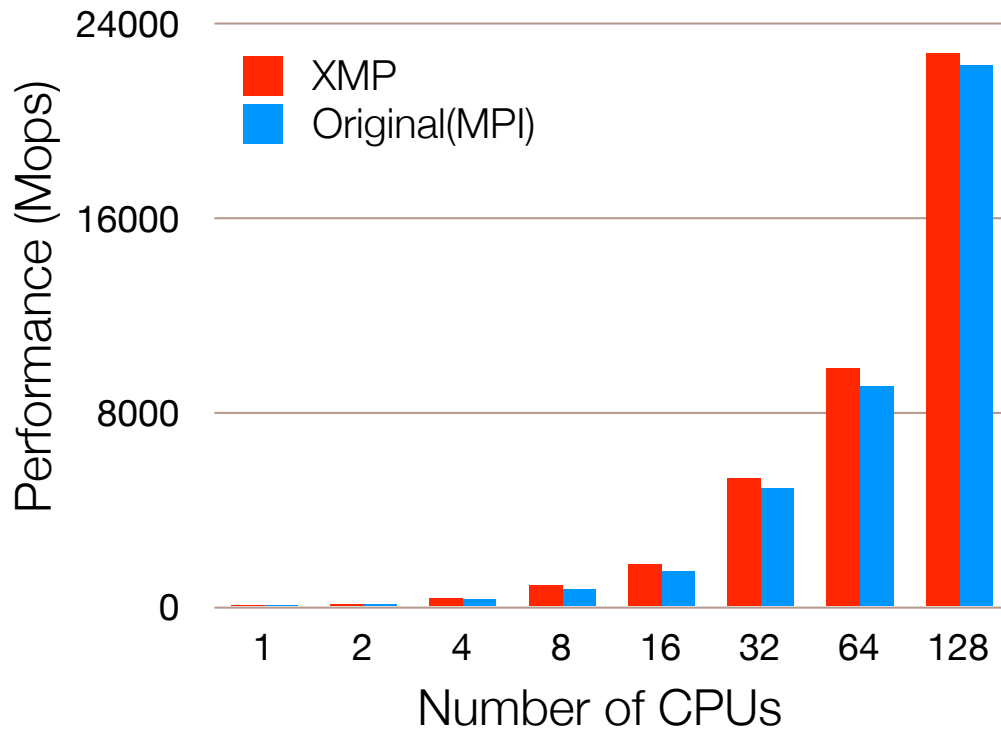
# Conjugate Gradient



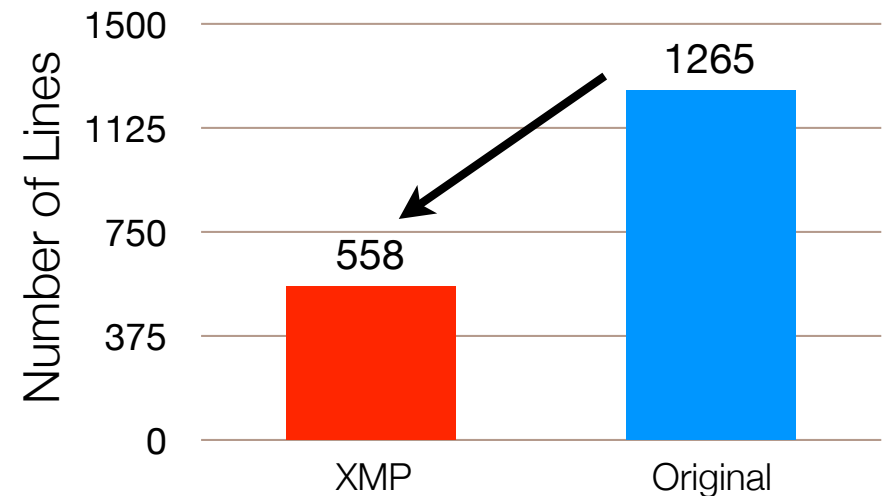
- Local view programming
- reduction directives for local variable

```
#pragma xmp coarray w, w1:[*]  
:  
for( i=ncols; i>=0; i-- ){  
    w[l:count[k][i]] += w1[m:count[k][i]][:k];  
}
```

## Performance



## Productivity





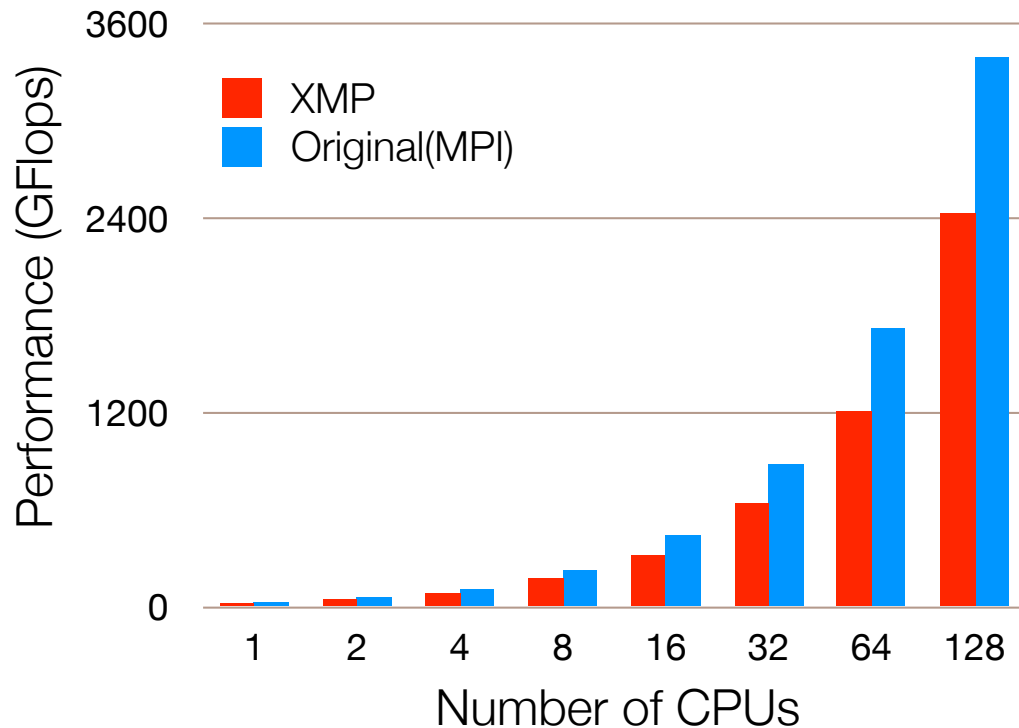
# High Performance Linpack



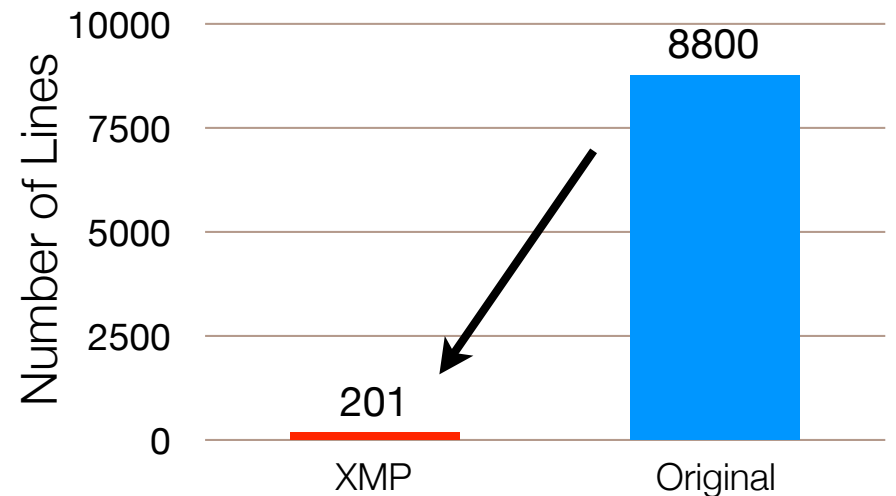
- block-cyclic distribution
- BLAS Lib. is used directly from distributed array

```
#pragma xmp distribute \  
    t(cyclic(NB), cyclic(NB)) onto p  
#pragma xmp align A[i][j] with t(j, i)  
:  
cblas_dgemm(..., &A[y][x], ...);
```

## Performance





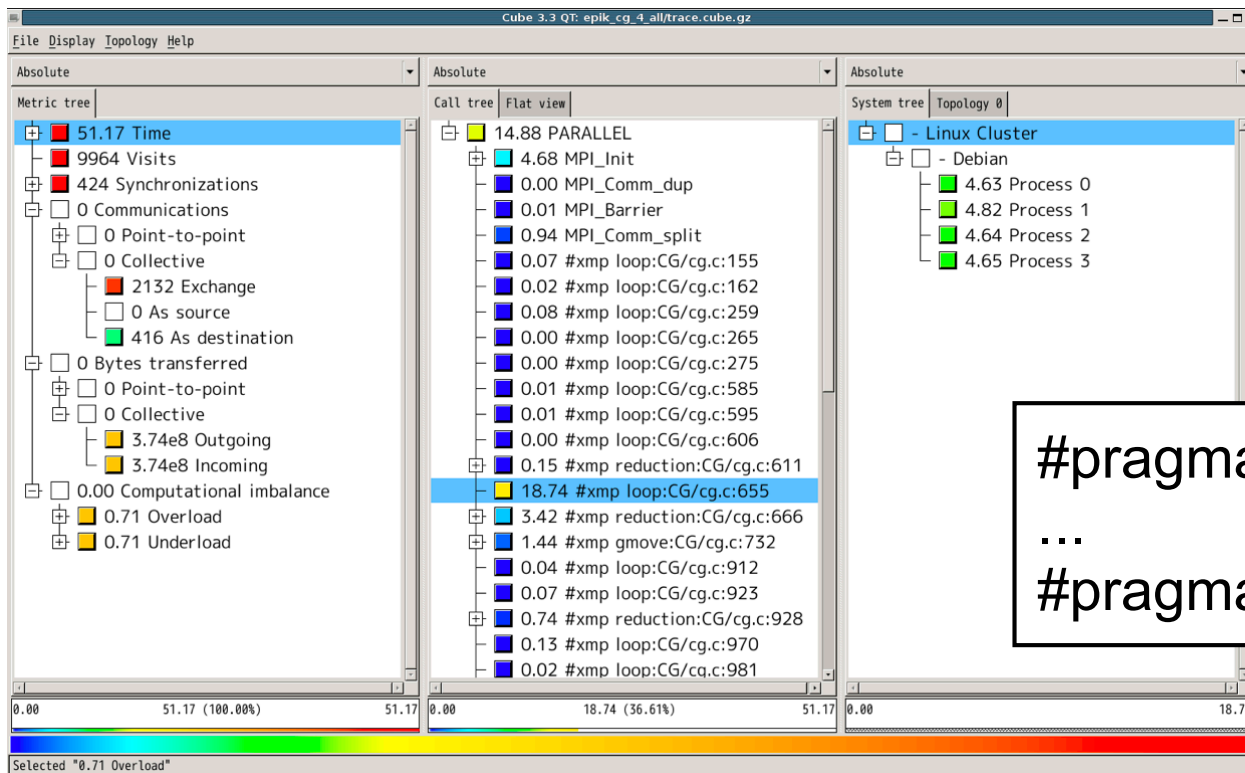
## Productivity



# International Collaboration



- Interface of XMP program profile to **Scalasca**
- **Scalasca** is a software tool that supports the performance optimization of parallel programs
- **Scalasca** is developed by  and 



```
#pragma xmp gmove profile
...
#pragma xmp loop on t(i) profile
```

# Summary & Future work



- XcalableMP was proposed as a new programming model to facilitate program parallel applications for distributed memory systems
- Evaluation of Performance and Productivity
  - Performance of XMP is compatible with that of MPI
  - Productivity of XMP is higher than that of MPI
- Future work
  - Performance evaluation on larger environment
  - For accelerators(GPU, etc), Parallel I/O , and Interface of MPI library