



2013 HPC Challenge Class 2 Submission

XcalableMP for Productivity and Performance

Masahiro Nakao[†], Hitoshi Murai[‡]
Takenori Shimosaka[‡], Mitsuhsa Sato^{†‡}

[†] Center for Computational Sciences, University of Tsukuba, Japan

[‡] RIKEN Advanced Institute for Computational Science, Japan



Outline



-
- 1.What's XcalableMP ?
 - 2.Implementations
 - 1. HPL
 - 2. RandomAccess
 - 3. FFT
 - 4. STREAM
 - 5. HIMENO Benchmark (Optional)
 - 3.Performances
 - 4.Conclusion

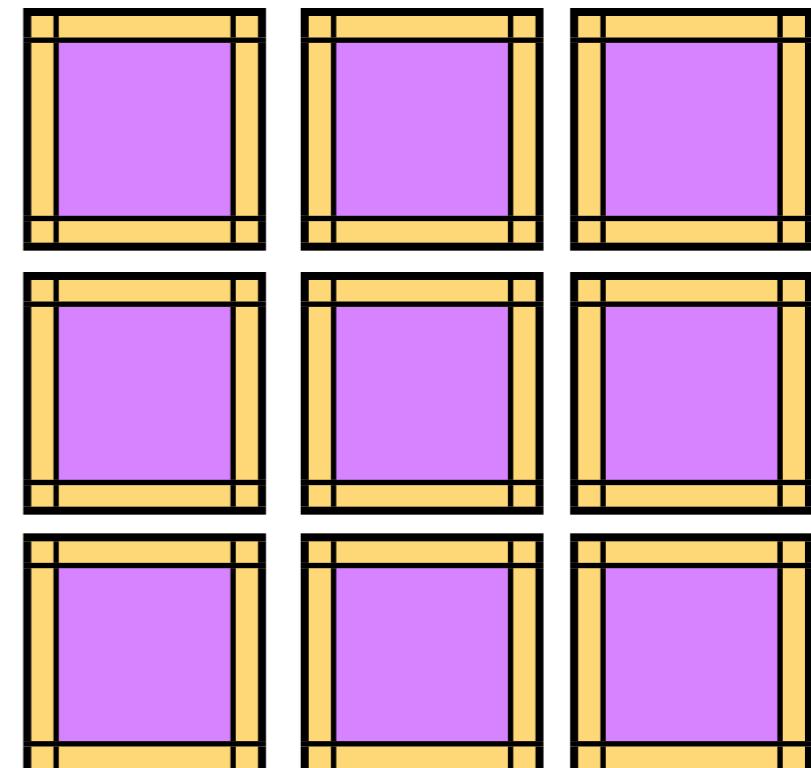
What is XcalableMP?

- XcalableMP (XMP)
 - Directive-based language extension of Fortran 2008 and C99
 - **XMP/Fortran** language and **XMP/C** language exist (same directives)
 - **Global-view** and **Local-view** memory models
 - XMP global-view model enables parallelizing the original sequential code using minimal modification with simple directives
 - **Coarray syntax** is available in both languages
 - XMP/Fortran is upward compatible with Fortran 2008
 - **XMP specification ver. 1.2** at <http://www.xcalablemp.org>
 - **XMP specification working group** consists of members from academia, research labs, and industries
 - Univ. of Tsukuba, RIKEN AICS, Fujitsu, NEC, Hitachi, and so on
 - **Omni XMP Compiler ver. 0.7** at <http://www.hpcs.cs.tsukuba.ac.jp/omni-compiler/xcalablemp/>

Benchmarks



- HPCC benchmarks : **HPL, RandomAccess, FFT, STREAM**
- Optional benchmark : **Himeno benchmark**
 - Typical stencil application (<http://accc.riken.jp/2444.htm>)
 - Evaluates performance of incompressible fluid analysis code
 - Why do we select the HIMENO benchmark ?
 - To demonstrate parallelization by XMP directives which synchronizes the overlapped regions



Benchmarks



Omni XcalableMP Compiler

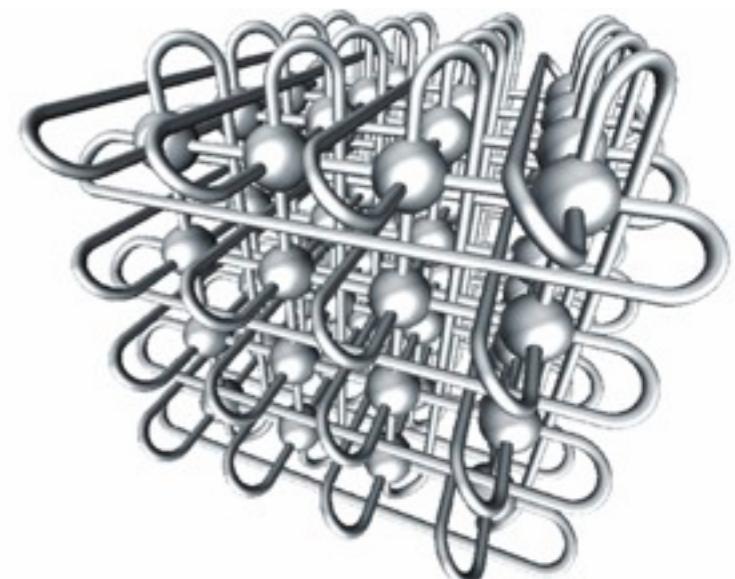
The screenshot shows a web browser window with the following details:

- Title Bar:** Omni XcalableMP Compiler
- Address Bar:** www.hpcs.cs.tsukuba.ac.jp/omni-compiler/xcalablemp/benchmark.html
- Left Sidebar (Menu):**
 - OpenMP Compiler **Omni**
 - HOME
 - Download
 - Document
 - Benchmark** (This option is highlighted with a red rounded rectangle)
 - Publication
 - Mailing List
 - Links
- Main Content Area:**
 - HPC Challenge Benchmarks**
 - [High Performance Linpack](#)
 - [RandomAccess](#)
 - [Fast Fourier Transform](#)
 - [EP STREAM triad](#)
 - HIMENO Benchmark**
 - [HIMENO Benchmark](#)
- Callout Box:** A black-bordered box contains the text: "You can download these implementations at <http://www.hpcs.cs.tsukuba.ac.jp/omni-compiler/xcalablemp/>"

The K computer



- SPARC64 VLIIfx Chip
 - **128 GFlops** (2GHz x 8Cores x 8)
 - 6MB L2 shared cache
 - 58W
- DDR3 SDRAM
 - **16GB**
 - 64GB/s
- Tofu Interconnect
 - 6D mesh/torus network
 - 5GB/s x 4links x 2
- **82,944** compute nodes (+ 5,184 IO nodes)



<http://www.aics.riken.jp>

Outline



1.What's XcalableMP ?

2.Implementations

- 1. HPL
- 2. RandomAccess
- 3. FFT
- 4. STREAM
- 5. HIMENO Benchmark (Optional)

3.Performances

4.Conclusion

Summary

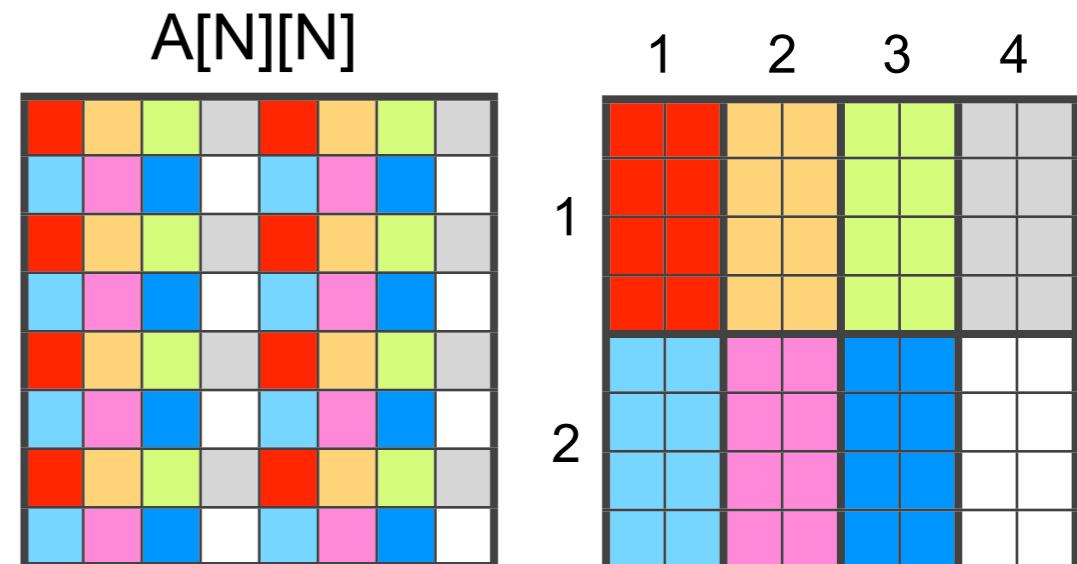
Benchmark	# Nodes	Performance	SLOC
HPL	16,384	933.8 TFlops (44.5% of peak)	306
RandomAccess	16,384	162.6 GUPs	250
FFT	36,864	50.1 TFlops (1.1% of peak)	239 + 283 + 1892
STREAM	16,384	481.8 TB/s	66
HIMENO	82,944	1.3 PFlops (12.7% of peak)	137



Full compute nodes

- Source lines of Code (SLOC) is **306**, written in XMP/C
- Block-Cyclic Distribution

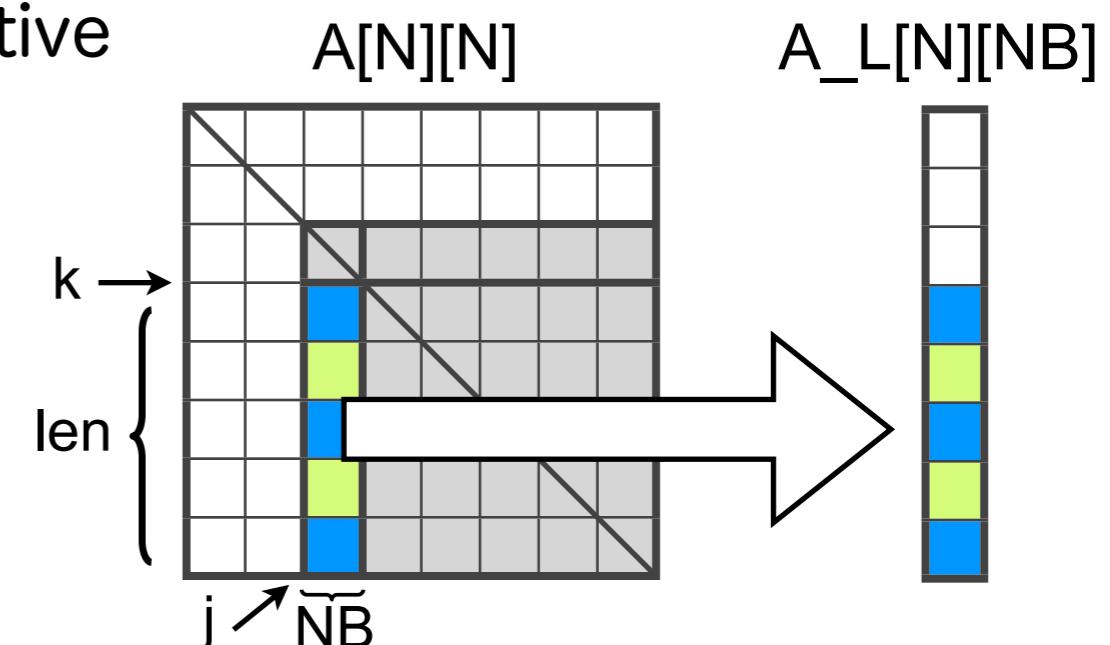
```
double A[N][N];
#pragma xmp nodes p(P,Q)
#pragma xmp template t(0:N-1, 0:N-1)
#pragma xmp distribute t(cyclic(NB), \
                        cyclic(NB)) onto p
#pragma xmp align A[i][j] with t(j,i)
```



Programmer can use BLAS for distributed array.

- Panel Broadcast by using **gmove** directive

```
double A_L[N][NB];
#pragma xmp align A_L[i][*] with t(*,i)
:
#pragma xmp gmove
A_L[k:len][0:NB] = A[k:len][j:NB];
```



RandomAccess



- SLOC is **253**, written in XMP/C
- The XMP RandomAccess is iterated over sets of CHUNK updates on each node
- Local-view programming with XMP/C coarray syntax

```
u64Int recv[LOGPROCS][RCHUNK+1]:[*];           ← Define coarray
...
for (j = 0; j < logNumProcs; j++) {
    recv[j][0:num]:[i_partner] = send[i][0:num]; ← Put operation
}
#pragma xmp sync memory
#pragma xmp post(p(i_partner), 0)
:
#pragma xmp wait(p(j_partner))
```

A diagram illustrating the XMP/C code structure. A vertical rectangle surrounds the entire code block. Three horizontal arrows point from the right side of the rectangle to specific parts of the code:

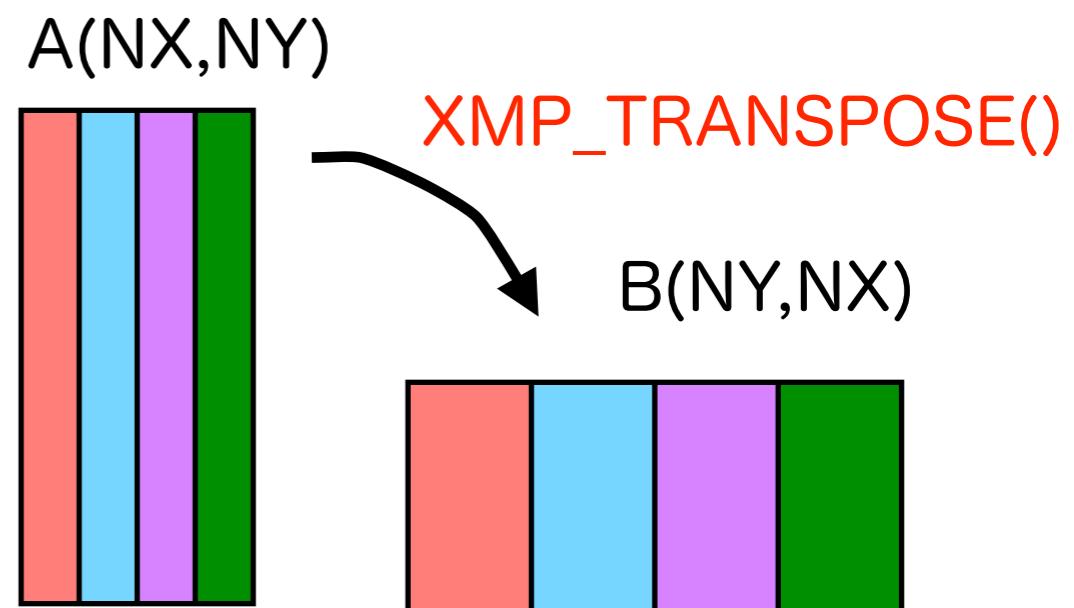
- An arrow points to the declaration `recv` with the label "Define coarray".
- An arrow points to the assignment statement `recv[j][0:num]:[i_partner] = send[i][0:num];` with the label "Put operation".
- An arrow points to the final `#pragma xmp wait` directive with an empty label.

- A point-to-point synchronization is specified with the XMP's `post` and `wait` directives to realize asynchronous behavior of this algorithm

- SLOC is **239**+ 283 +1,892, written in **XMP/Fortran + OpenMP**
 - We mainly modified “pzfft1d.f” in ffte optimized for the K computer
 - This SLOC is **70** (the SLOC of original “pzfft1d.f” is 101)
 - The total SLOC of the modified files is 239 (**70** + 169)
 - The total SLOC of NOT modified files is 283
 - The total SLOC of C interface and another kernels of the ffte is 1,892

```
CALL XMP_TRANSPOSE(B,A,1)
...
!$XMP loop on tx(I)
!$OMP parallel do
DO I=1,NX
  DO J=1,NY
    B(J,I)=B(J,I)*W(J,I)
  END DO
END DO
```

A() and B() are global arrays



EP STREAM Triad



- SLOC is **66**, written in **XMP/C + OpenMP**
- a[], b[], and c[] are local arrays

```
#pragma xmp nodes p(*)
...
for(k=0; k<NTIMES; k++) {
    #pragma omp parallel for
    for (j=0; j<size; j++)
        a[j] = b[j] + scalar*c[j];
}
...
*triadGBs = ... // performance of each node
...
#pragma xmp reduction(+:triadGBs)
```

Parallelize the program

Gather performance results
on each node

- SLOC is **137**, written in XMP/Fortran
 - SLOC of the original is **380**, written in MPI/Fortran

To synchronize overlapped regions, we use XMP shadow and reflect directives

```

!$xmp shadow p(0,2:1,2:1)
...
!$xmp reflect (p)
...
!$xmp loop (J,K) on t(*,J,K)
do K = 2, kmax-1
    do J = 2, jmax-1
        do I = 2, imax-1
            S0 = a(I,J,K,1)*p(I+1,J,K) + ...
            :
        enddo
    enddo
enddo

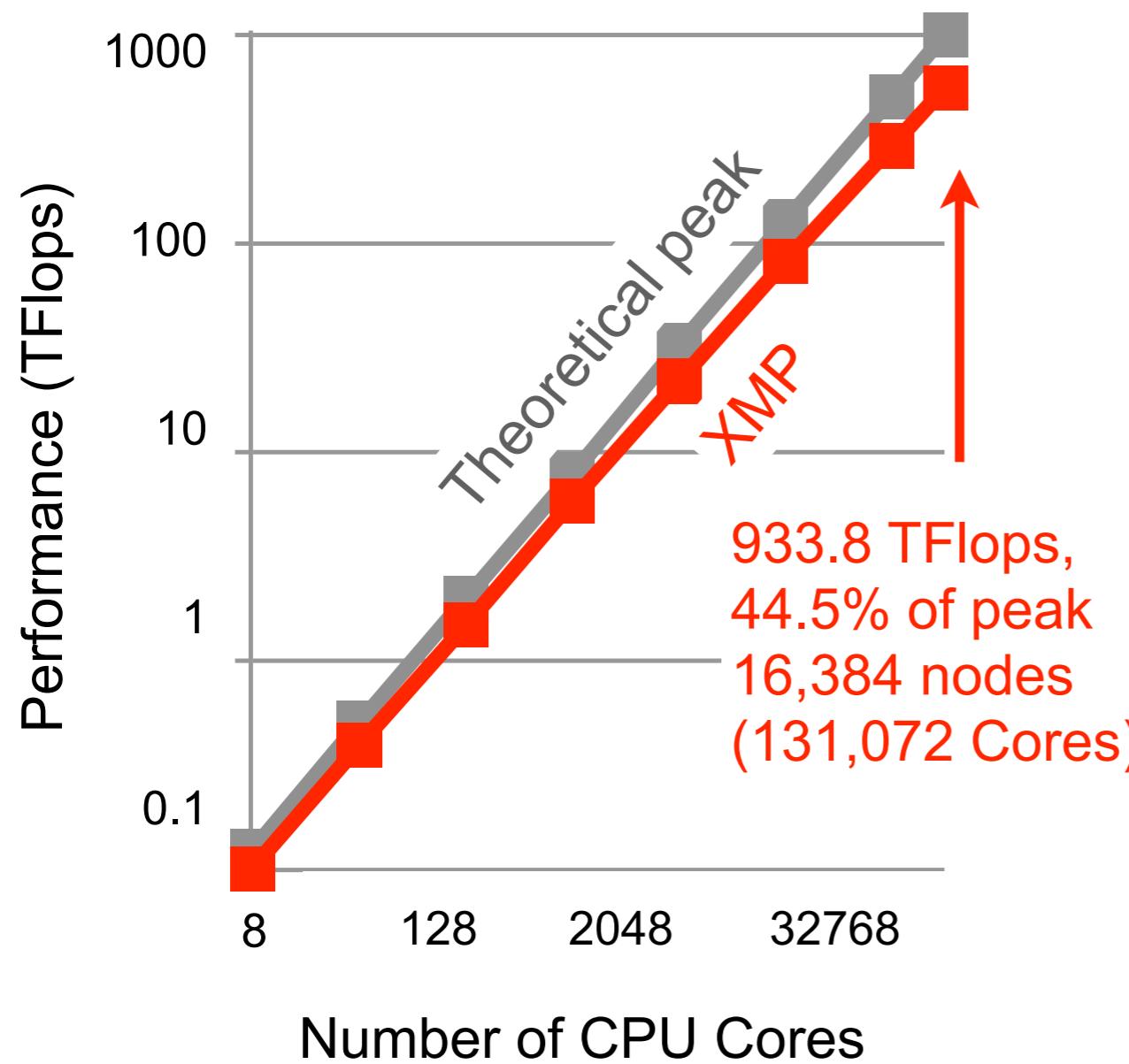
```

Define overlapped region
 Synchronize overlapped region
 Work mapping

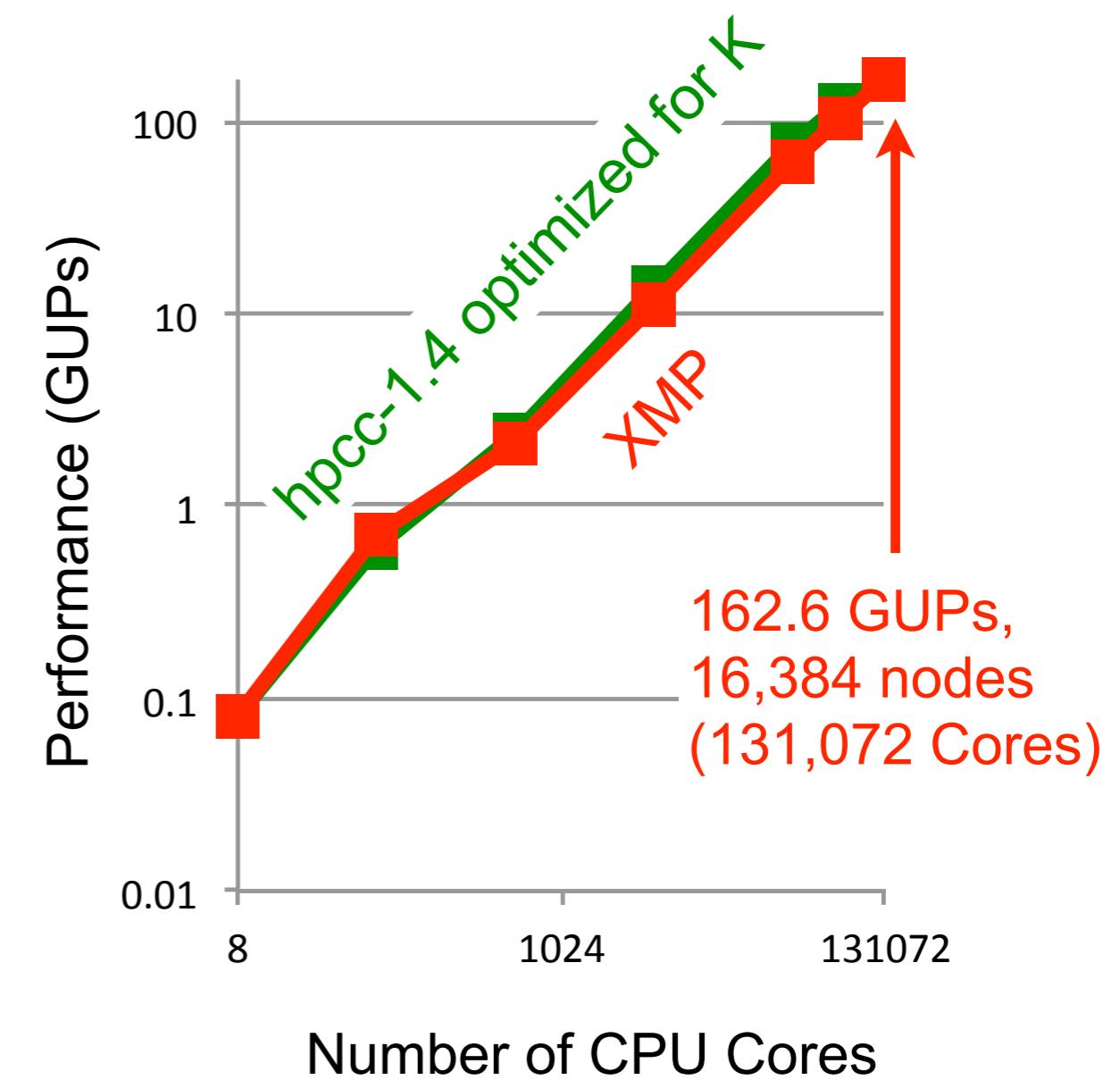
Only add XMP directives into the sequential Himeno benchmark basically.

Results(1/3)

- HPL (1 process/node + Threaded BLAS)



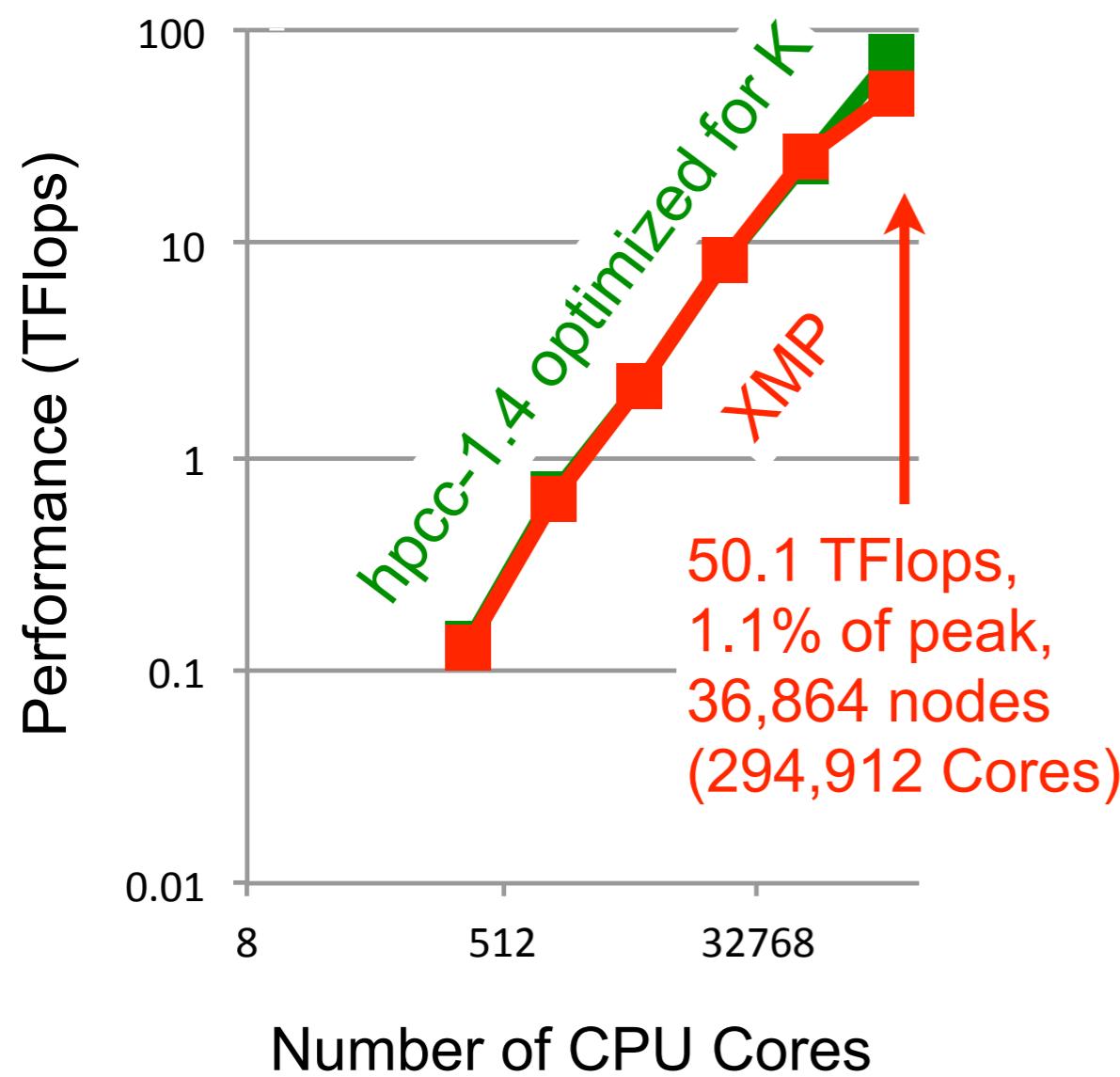
- RandomAccess (8 processes/node)



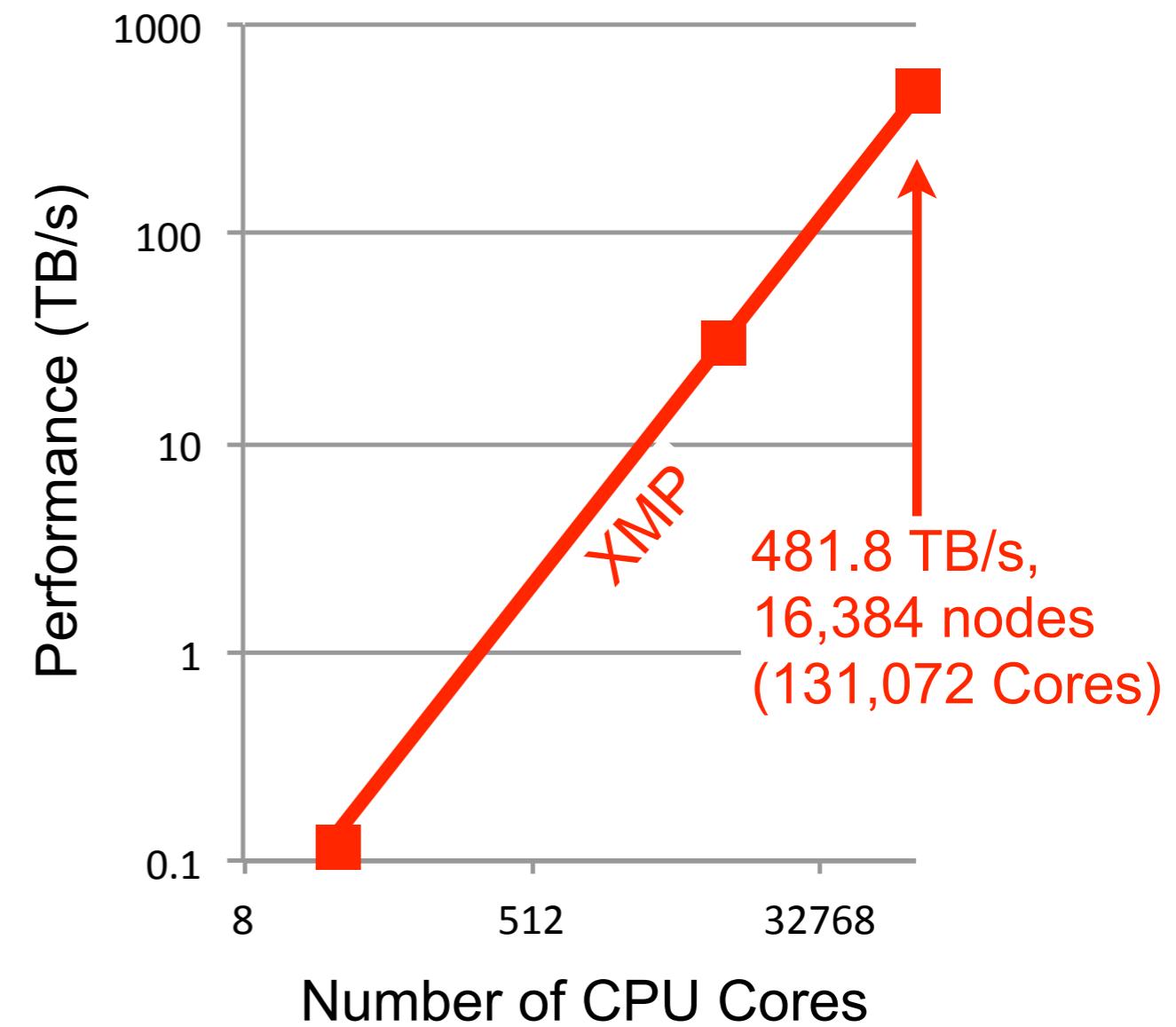
Results(2/3)



- FFT (1 process/node with 8 threads)



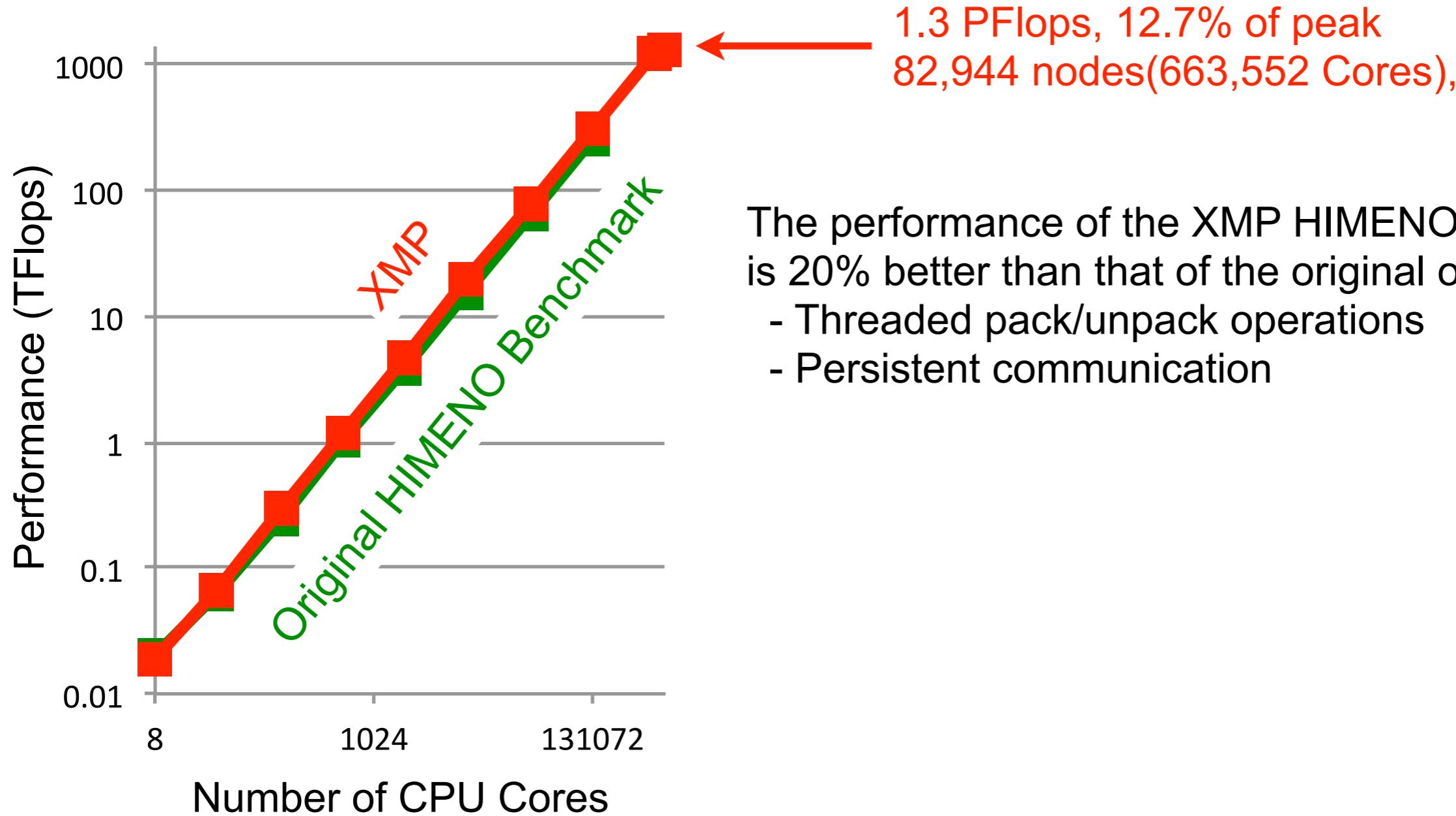
- STREAM (1 process/node with 8 threads)



Results(3/3)



- HIMENO Benchmark
(1 process/node with 8 threads)



The performance of the XMP HIMENO Benchmark is 20% better than that of the original one.
- Threaded pack/unpack operations
- Persistent communication

Conclusion



Benchmark	# Nodes	Performance	SLOC
HPL	16,384	933.8 TFlops (44.5% of peak)	306
RandomAccess	16,384	162.6 GUPs	250
FFT	36,864	50.1 TFlops (1.1% of peak)	239 + 283 + 1892
STREAM	16,384	481.8 TB/s	66
HIMENO	82,944	1.3 PFlops (12.7% of peak)	137

Masahiro Nakao, et al. "Productivity and Performance of the HPC Challenge Benchmarks with the XcalableMP PGAS language", 7th International Conference on PGAS Programming Models, October, 2013

For more information, please visit

- #1437 RIKEN AICS (Advanced Institute for Computational Science)**
- #2519 Center for Computational Sciences, University of Tsukuba**